

# Confidentiality and Sequencing in Bilateral Negotiations

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Many bargaining problems involve one buyer and several sellers:

Real Est. Dev.	→	Land Owners
Econ. Dept.	→	Multiple Faculty
Home-owner	→	Multiple Contractors

- In many applications, the buyer can choose between public and private negotiations.
  - Disney was very secret about purchases of Florida swamp land.
  - Duke was not as lucky in a small community.
- For government buyers, “sunshine” laws require transparency.

*“[These laws] are the state and federal statutes requiring that government meetings, decisions and records be made available to the public.”*

- And sunshine laws have a bite.
- For instance, in July 2000, Motorola alleged that the state of Florida violated its own sunshine law by holding closed meetings with the rival Com-Net.

- Exceptions, however, apply. For instance, during real property acquisitions by the state's water management districts,

*“...appraisal reports, offers, and counteroffers are confidential and exempt from the provisions of [sunshine law] until an option contract is executed.” (Florida Statutes 373.139)*

- The buyer can also choose the sequence of negotiations.
- For example, is it better to start with
  - the weaker or stronger seller.
  - the more or less valuable item, etc.?

- **Confidential Negotiations**

Calzolari and Pavan (*JET*, 2006); Krasteva and Yildirim (*GEB*, 2012); Noe and Wang (*ReStud*, 2004); Taylor (*RAND*, 2004).

- **Buyer's Sequencing Preference**

Cai (*JET*, 2000); Marx and Shaffer (*IJIO*, 2007); Krasteva and Yildirim (*RAND*, Forth.); Li (2010); Marshall and Merlo (*IER*, 2004); Moresi, Salop, and Sarafidis (2010); Xiao (2010).

- **Seller's Sequencing Preference**

Arbatskaya (*RAND*, 2007); Marx and Shaffer (*IJIO*, 2010).

# Model (w/ Perfect Complements)

Based on Krasteva and Yildirim (GEB, 2012)

- Three risk-neutral players: one buyer ( $b$ ) and two sellers ( $s_i$ ,  $i = 1, 2$ ).
- Buyer's valuations are
  - $V = 1$  from a joint purchase,
  - $v_1 = v_2 = 0$  from each unit.
- $c_i = 0 \implies$  *efficient* to purchase both.
- Bargaining over price,  $p_i$ :

$$\begin{cases} s_i, & \text{w/ prob. } \alpha \\ b, & \text{w/ prob. } 1 - \alpha. \end{cases}$$



# Public vs. Private Negotiations

eXploding offers (i.e., pay-as-you-go)

- **Public.**  $p_1(s_1) = 1 - \alpha$  and  $p_2(s_2) = 1 \implies$  efficient trade.

And, the buyer's expected payoff is:

$$\begin{aligned}\pi^X(b) &= (1 - \alpha)^2 + \alpha(1 - \alpha)[1 - (1 - \alpha + 0)] \\ &\quad + (1 - \alpha)\alpha[1 - (0 + 1)] + \alpha^2[1 - (1 - \alpha + 1)] \\ &= (1 - \alpha)^2.\end{aligned}$$

- **Private.**  $p_1(s_1) = p_2(s_2) = \frac{1}{1+\alpha}$ , and the buyer visits each seller w/ prob.  $\frac{1}{2}$ .

$\implies$  efficient trade. And,

$$\bar{\pi}^X(b) = 1 - 2\alpha \frac{1}{1 + \alpha} = \frac{1 - \alpha}{1 + \alpha} > (1 - \alpha)^2 = \pi^X(b).$$

# Public vs. Private Negotiations

Open-ended offers (i.e., decision at the end)

- **Public.**  $p_1(s_1) = 1$  and  $p_2(s_2) = 0 \implies$  efficient trade.

And,

$$\pi^O(b) = (1 - \alpha)^2.$$

- Assume  $\alpha < \frac{1}{2}$ .
- **Private.**  $p_1(s_1) = p_2(s_2) = 1 \implies$  inefficient trade.

And,

$$\bar{\pi}^O(b) = (1 - \alpha)^2.$$

- Public neg. is efficient while private neg. can be inefficient.
- The buyer (weakly) prefers private neg.
- May explain the presence of “sunshine” laws.
- Extended return and cancellation policies, e.g., FTC’s cooling-off rule may hinder efficiency.

# Sellers' Incentives to Disclose

- Consider *eXploding* offers, and suppose sellers pre-commit to their disclosure policies.



$s_i \backslash s_j$	$dis$	$\overline{dis}$
$dis$	$\alpha(1 - \frac{\alpha}{2}); \alpha(1 - \frac{\alpha}{2})$	$\alpha; \alpha(1 - \alpha)$
$\overline{dis}$	$\alpha(1 - \alpha); \alpha$	$\alpha \frac{1}{1+\alpha}; \alpha \frac{1}{1+\alpha}$

- Unique NE =  $(dis, dis) \implies$  public neg.

# Strategic Sequencing

- Consider the sequence  $s_i \rightarrow s_j$ .
- Let  $v_1 = v_2 = v < \frac{1}{2}$ .
- Assume fully contingent price contracts.
- Then, the buyer and  $s_i$  can always conspire to leave 0 surplus to  $s_j$ .
- But how?

**Proposition.** Given the sequence  $s_i \rightarrow s_j$ , the following prices constitute a SPNE:

- On-equilibrium path:
  - efficient trade; 0 profit for  $s_j$  (similar to Aghion and Bolton (*AER*, 1987)).

$$\begin{aligned}p_i^{ij}(s_i) &= 1 - (1 - \alpha_j)v \text{ and } p_j^{ij}(s_j) = 0 \\p_i^{ij}(b) &= p_j^{ij}(b) = 0.\end{aligned}$$

- Off-equilibrium path:
  - likely below-cost pricing and “breakup” fee.

$$p_i^i(b) = -(1 - v) \text{ and } p_i^i(s_i) = p_i^{-i}(s_i) = \alpha_j v.$$

- 1 The buyer receives expected payoff:

$$\pi_{ij}(b) = (1 - \alpha_i) + \alpha_i(1 - \alpha_j)v.$$

- 2  $\pi_{ij}(b) > \pi_{ji}(b) \iff \alpha_i < \alpha_j.$

$\implies$  start with the *weaker* seller.

# Non-contingent Prices

- The unique (equilibrium) prices are:

$$p_i^*(s_i) = 1 - v \text{ and } p_j^*(s_j|1 - v) = v.$$

- The buyer's payoff is:

$$\pi_{ij}(b) = (1 - \alpha_i)(1 - \alpha_j) + \alpha_i(1 - \alpha_j)v + (1 - \alpha_i)\alpha_jv.$$

- Indifference to the sequence!



# Model (w/ Payoff Uncertainty)

Based on Krasteva and Yildirim (RAND, Forth.).

- $V = 1$  from a joint purchase.
- $v_i = \begin{cases} 0, & \text{w/ prob. } q \\ \frac{1}{2}, & \text{w/ prob. } 1 - q. \end{cases}$

**Proposition.** Fix the sequence  $s_i \rightarrow s_j$ . In equilibrium, if  $s_i$  makes the offer for product  $i$ , then

$$(p_i^*(s_i), p_j^*(s_j|s_i)) = \begin{cases} (1, \frac{1}{2}) & \text{if } \alpha_j < \hat{\alpha}(q) \\ \left(\frac{1+q}{2}, \frac{1-q}{2}\right) & \text{if } \alpha_j \geq \hat{\alpha}(q) \text{ and } q > \frac{\sqrt{5}-1}{2} \\ \left(\frac{1}{2}, \frac{1}{2}\right) & \text{if } \alpha_j \geq \hat{\alpha}(q) \text{ and } q \leq \frac{\sqrt{5}-1}{2}. \end{cases}$$

**Proposition.** *Equilibrium trade is efficient irrespective of the sequence if and only if  $q \notin (\frac{1}{2}, 1)$ , or equivalently  $\hat{\alpha}(q) = 0$ . In addition, if  $q \notin (\frac{1}{2}, 1)$ , then the buyer is indifferent to the sequencing.*

**Proposition.** Let  $q \in (\frac{1}{2}, 1)$  and  $\alpha_1 < \alpha_2$ . Then,

$$\left\{ \begin{array}{ll} \text{is indifferent to the sequence} & \text{if } \alpha_1 < \alpha_2 < \hat{\alpha}(q) \\ \text{strictly prefers the sequence } s_1 \rightarrow s_2 & \text{if } \alpha_1 < \hat{\alpha}(q) \leq \alpha_2 \\ \text{strictly prefers the sequence } s_2 \rightarrow s_1 & \text{if } \hat{\alpha}(q) \leq \alpha_1 < \alpha_2. \end{array} \right.$$

# Raising Own Cost of Acquisition

## Minimum Purchase Price

**Proposition.** *Let  $q_1 = q_2 = q$  and  $\alpha_1 = \alpha_2 = \alpha$ . Suppose that prior to negotiations, the buyer commits to pay at least  $w \geq 0$  for each unit she purchases. Then,*

- *for  $q \in (\frac{1}{2}, 1)$ , there exists  $\alpha^c(q) \in (0, \hat{\alpha}(q))$  such that the buyer optimally sets  $w > 0$  for all  $\alpha \in [\alpha^c(q), \hat{\alpha}(q))$ ;*
- *for  $q \notin (\frac{1}{2}, 1)$ , the buyer optimally sets  $w = 0$  for all  $\alpha$ .*

# Raising Own Cost of Acquisition

## Strategic Outsourcing

**Proposition.** *Let  $\alpha_1 = \alpha_2 = \alpha$ , and suppose that the buyer can costlessly make input 1 in-house, which has no stand-alone value,  $q_1 = 1$ , but she needs to outsource input 2 with  $q_2 \in (\frac{\sqrt{5}-1}{2}, 1)$ . Then, the buyer is strictly better off outsourcing both inputs than only input 2 if and only if*

$$\alpha > \frac{1+q_2^2}{1+q_2}.$$

## Two sets of observations.

- ① Efficiency vs. Sequencing
  - Efficient trade implies indifference to sequencing
  - Under Inefficient trade, the sequence is important.
- ② Raising own cost of acquisition
  - Setting a minimum purchase price
  - Outsourcing even when available in-house.